


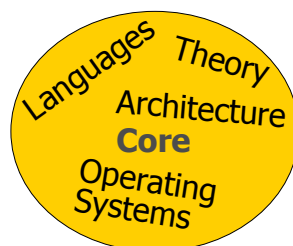
Real-Time Sensor Networks: Paradigms, Challenges, and Open Issues

Tarek Abdelzaher
*Department of Computer Science
University of Virginia*

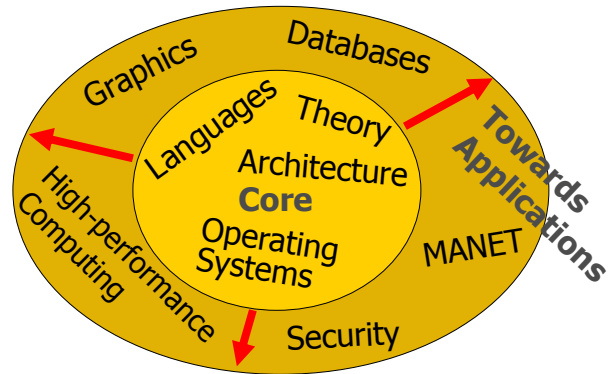


Where is Computer Science Research Going?

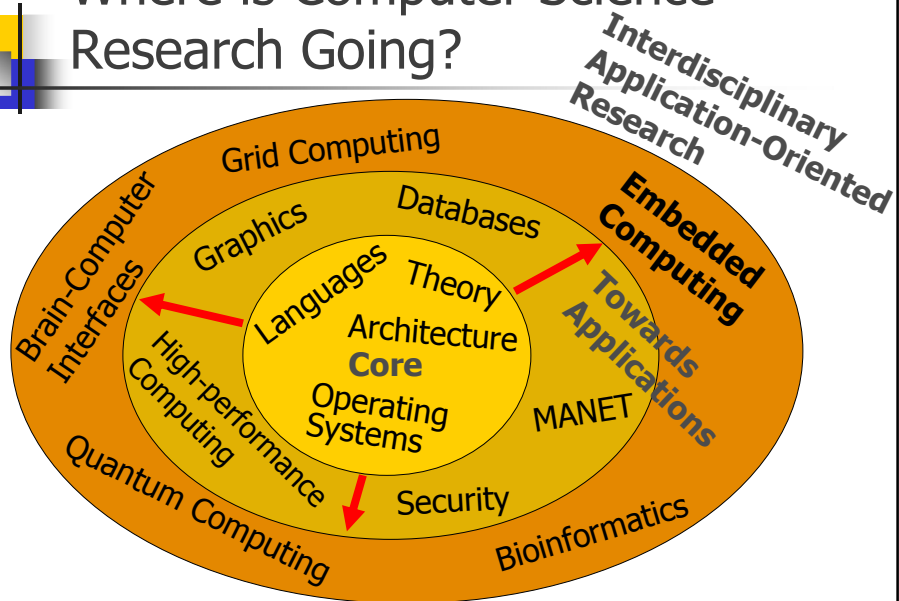
The beginning:



Where is Computer Science Research Going?



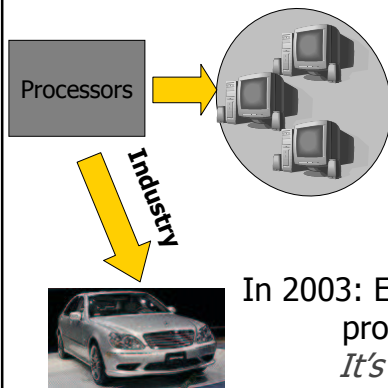
Where is Computer Science Research Going?



Real-Time and Embedded Systems

The Next Frontier

Why embedded computing?



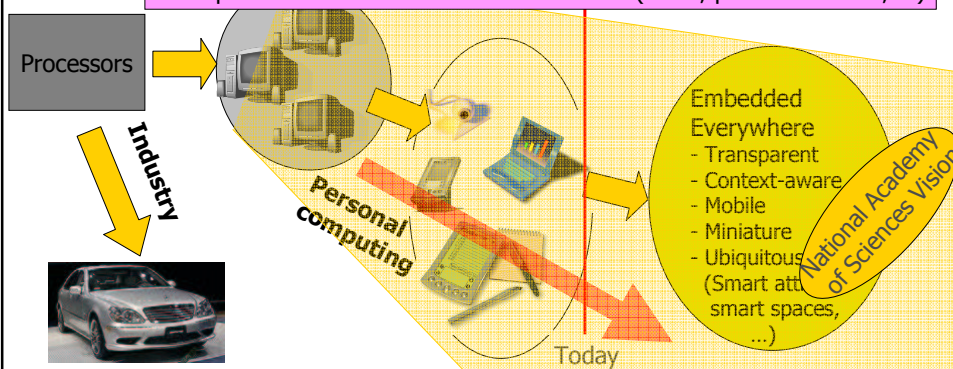
In 2003: Embedded processors are 98% of all processors manufactured
It's only the beginning...

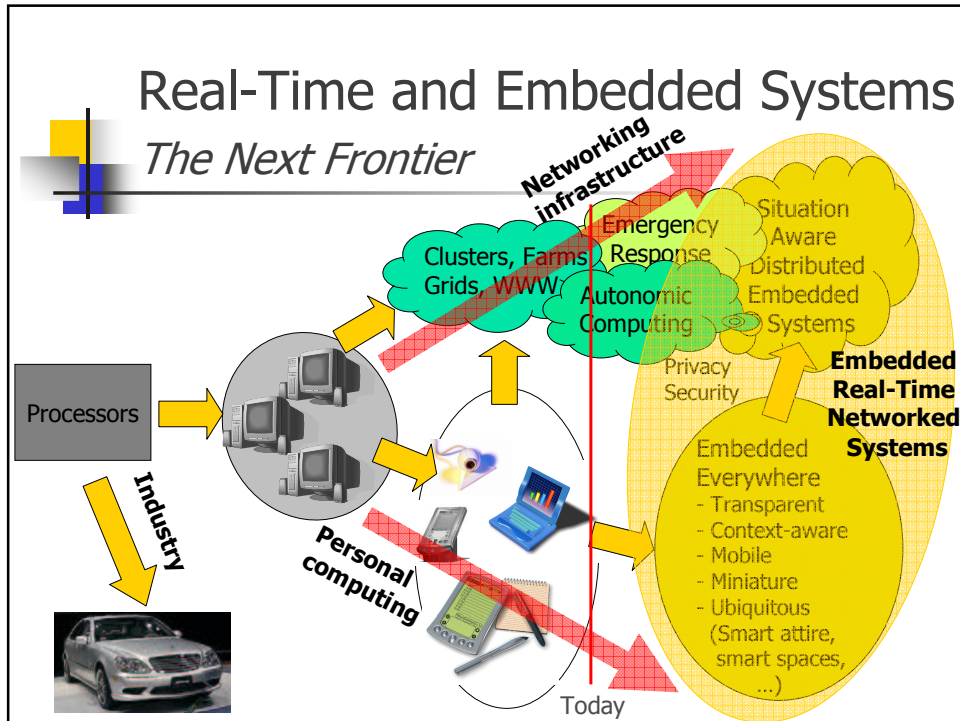
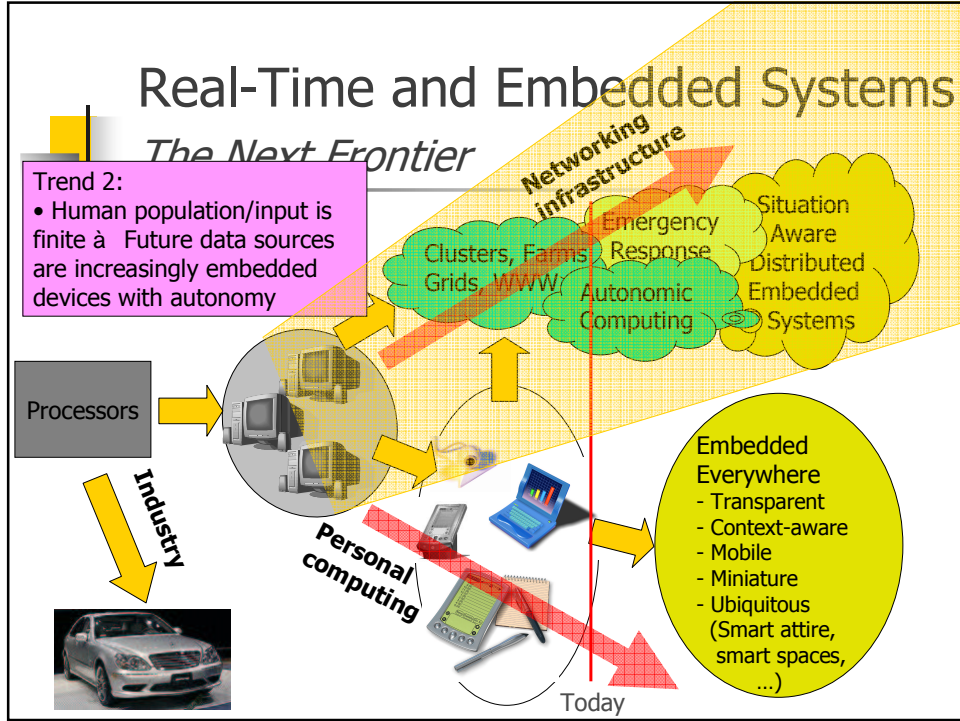
Real-Time and Embedded Systems

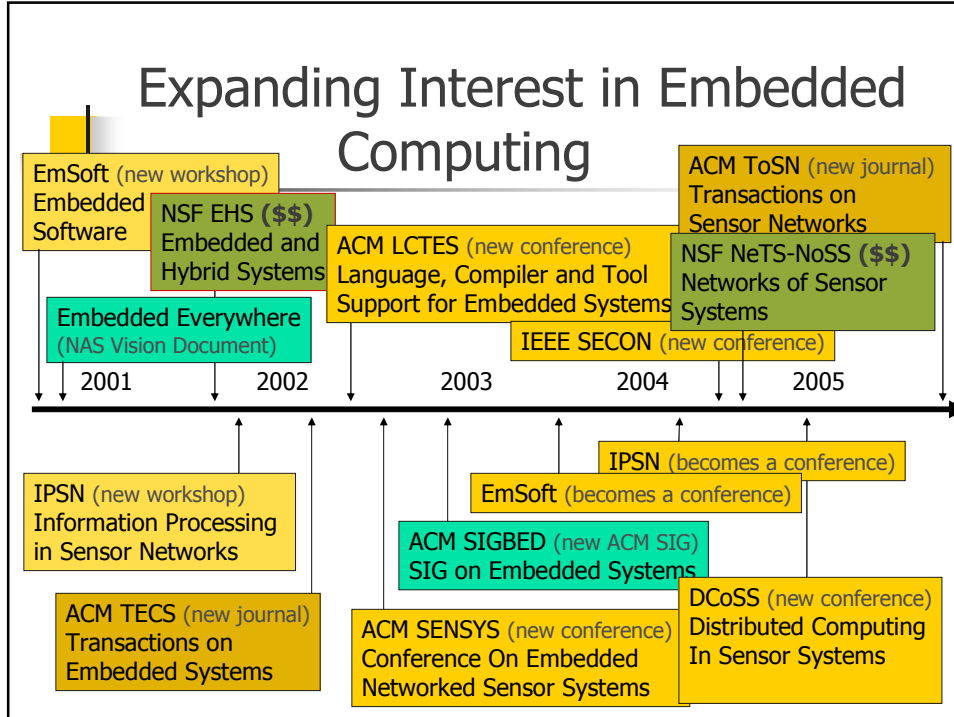
The Next Frontier

Trend 1:

- Invisible (embedded) computing, implicit interfaces (users need only 1 mobile device – rest should be non-intrusive)
- Context-aware computing (new sensors, new effectors)
- Ubiquitous – instrument what we use most (attire, personal effects, ...)







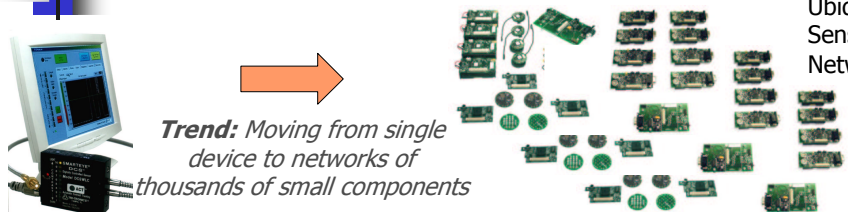
Overarching Challenge: A Science of *Design and Performance Analysis* for Embedded Networks

- n An NSF initiative: Science of Design
 - n "How can we explain, ***predict and control emergent properties*** of software-intensive systems?"
 - n "To what extent can we systematize guidance that leads to systems that ***satisfy requirements?***"
 - n "How can we develop theories that rely on aggregate ***reasoning about overall behavior*** rather than exact reasoning about all the details?"
- n A sensor networks perspective
 - n A science of design ***within constraints of time, space, and laws-of-physics***

Sensor Network Design

Three Software Challenges

Vision:
Ubiquitous
Sensor
Networks



Software Challenges:

Programming
Paradigms and
Protocols

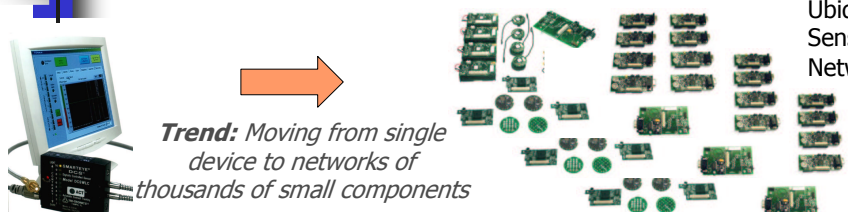
Analysis of aggregate
properties (lifetime,
capacity, ...)

Controllable emergent behavior

Sensor Network Design

Three Software Challenges

Vision:
Ubiquitous
Sensor
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Software Challenges:

Programming
Paradigms and
Protocols


Analysis of aggregate
properties (lifetime,
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Controllable emergent behavior

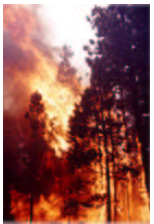
Application Domain




Habitat Monitoring




Precision Agriculture



Disaster Response




Target Tracking



Infrastructure Protection

Features

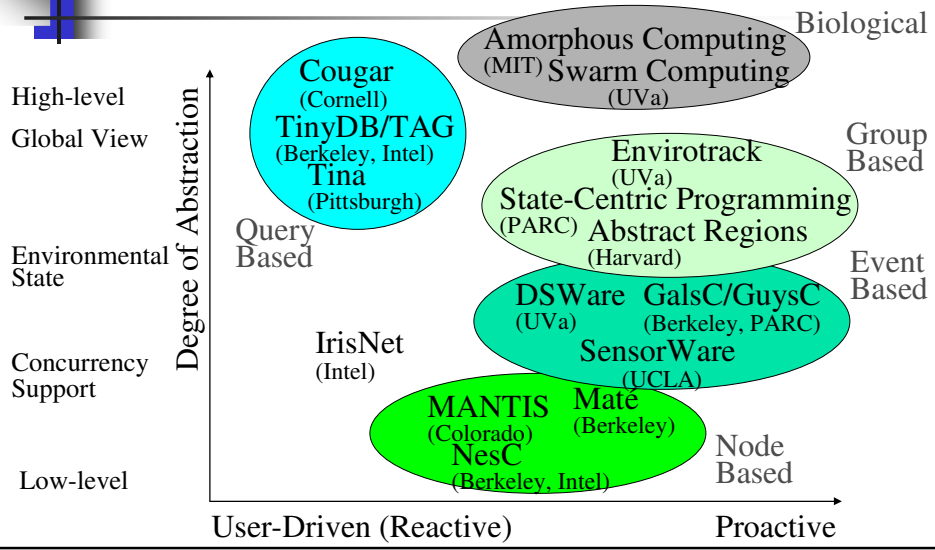
- Ad hoc deployment
- Massive distribution
- Interaction with a physical environment
- Unattended operation



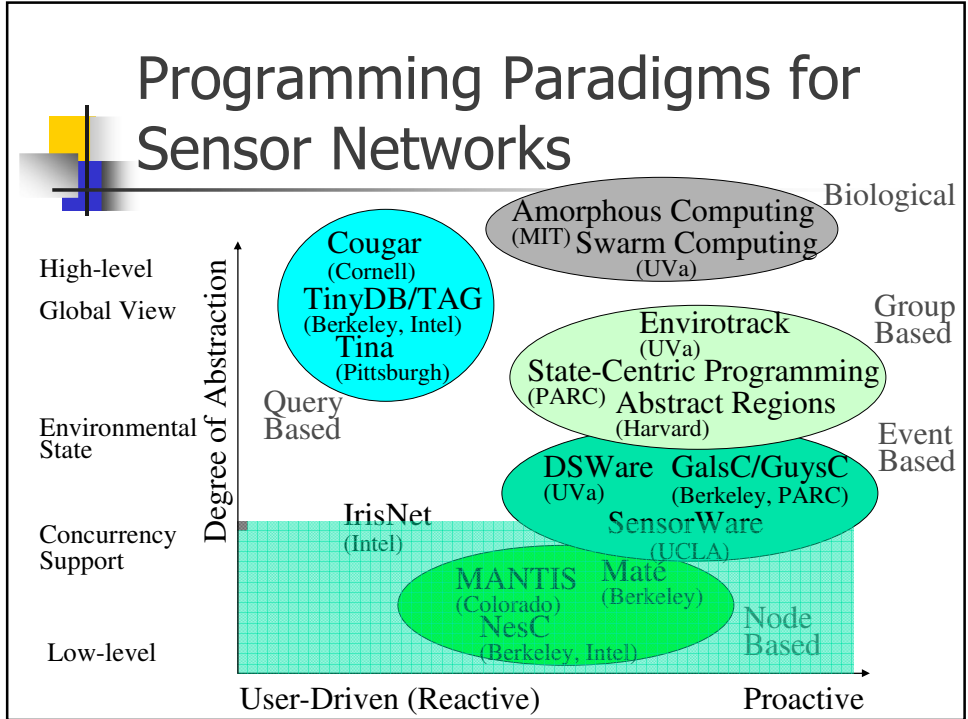
American Border Patrol

Border Control

Programming Paradigms for Sensor Networks

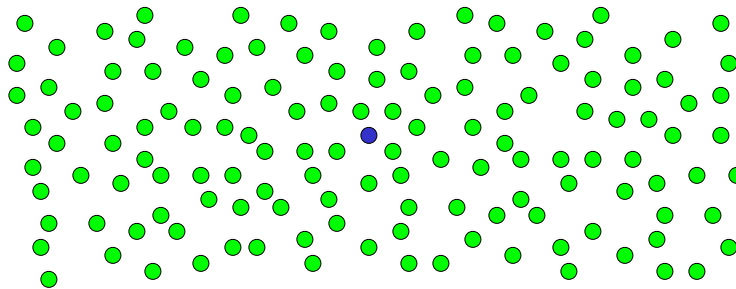


Programming Paradigms for Sensor Networks



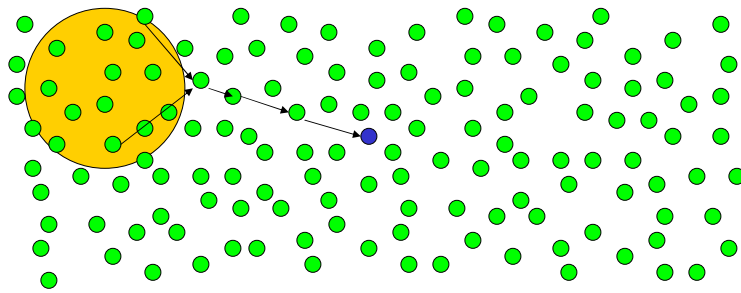
Challenge 1: Energy Balancing

n Problem: data collection by centralized entities



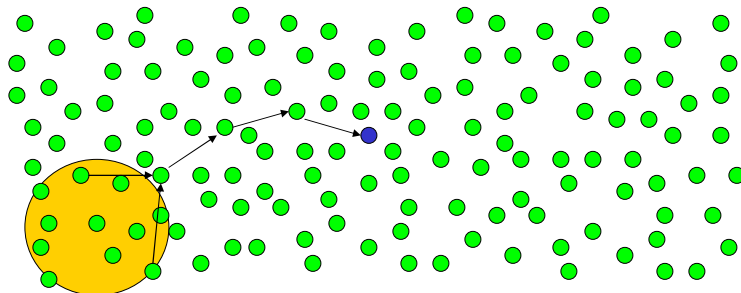
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n Problem: data collection by centralized entities



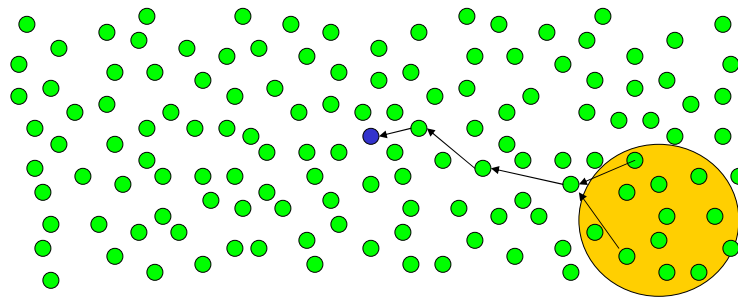
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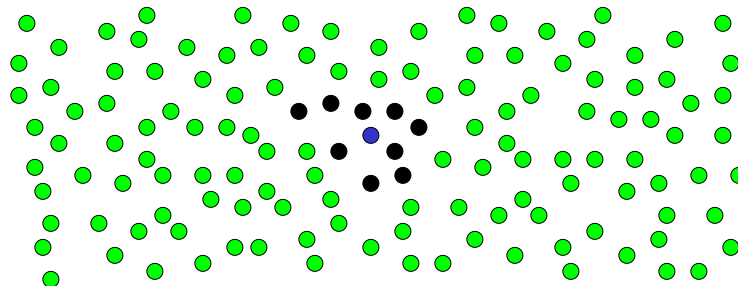
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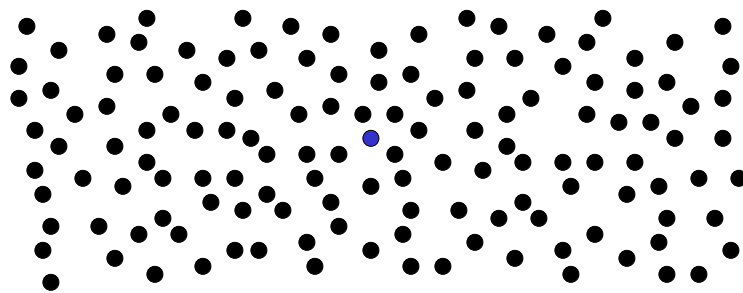
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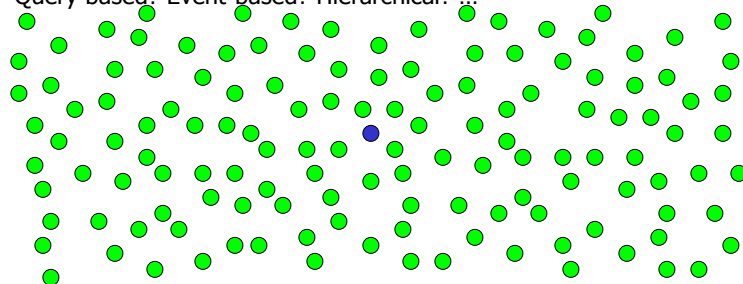
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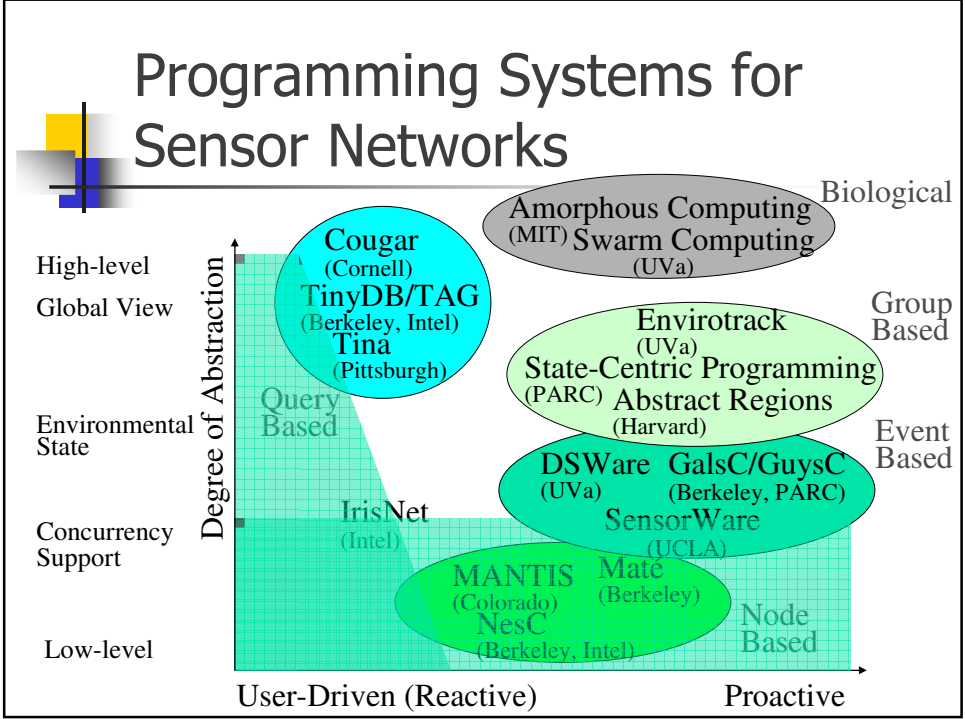
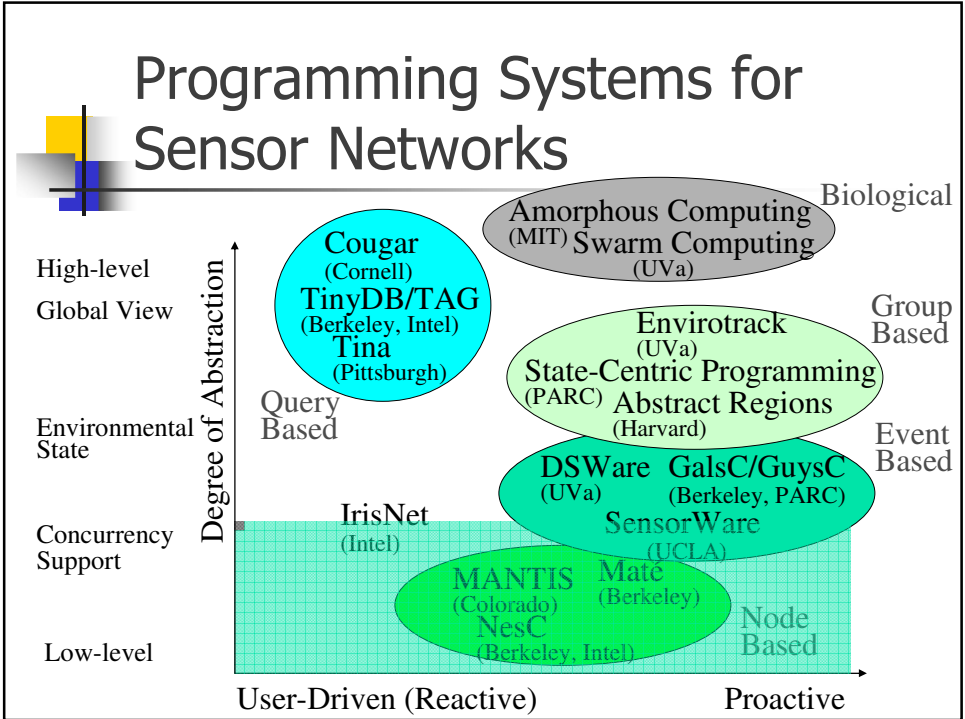
- n Problem: data collection by centralized entities



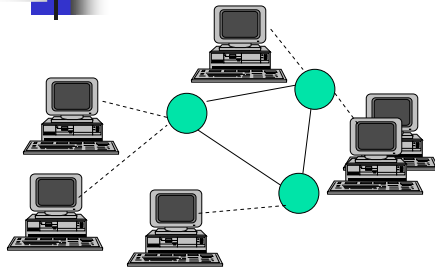
Challenge 1: Energy Balancing

- n Problem: data collection by centralized entities
- n Implications:
 - n Programming paradigms should support autonomous operation
 - n Push computation into the network
 - n Take the user out of the loop
 - n Query based? Event-based? Hierarchical? ...

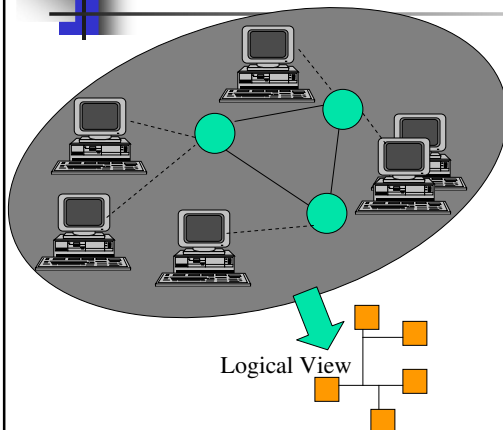




Challenge 2: Environmental Abstractions

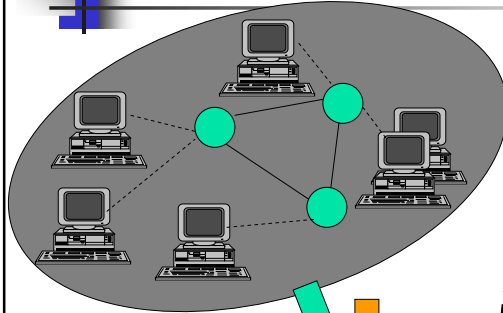


Challenge 2: Environmental Abstractions



- n Distributed programming paradigms
 - n Abstract distributed communication
 - n Provide location transparency

Challenge 2: Environmental Abstractions



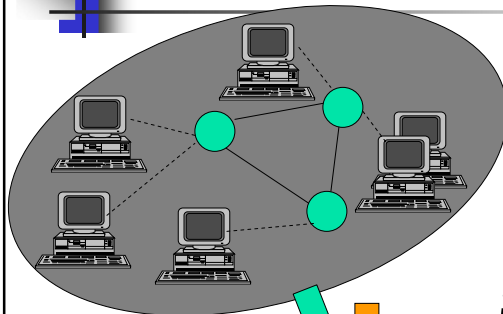
Logical View

- Distributed programming paradigms
 - Abstract distributed communication
 - Provide location transparency



- Sensor Network Programming Abstractions
 - Represent the physical world to the programmer
 - Abstract distributed interaction with the physical environment

Challenge 2: Environmental Abstractions

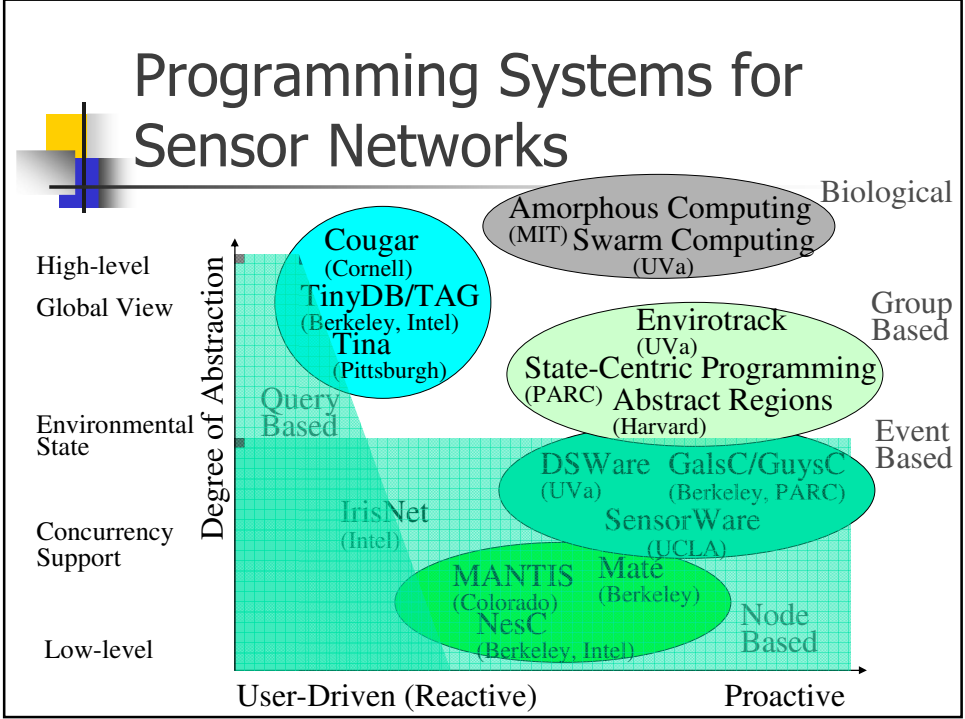
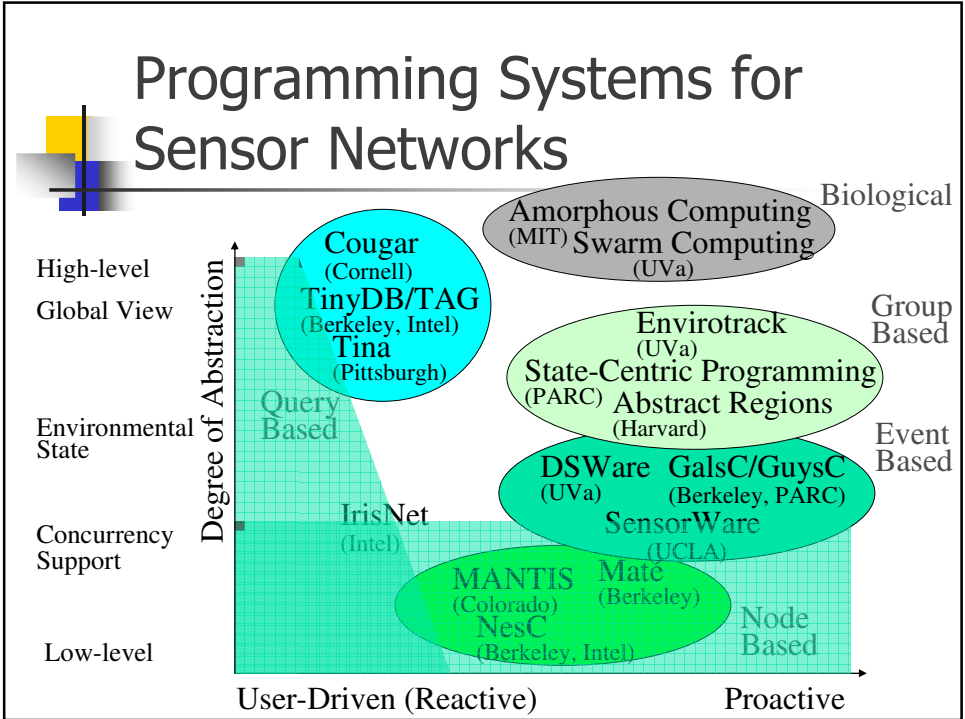


Logical View

- Distributed programming paradigms
 - Abstract distributed communication
 - Provide location transparency




- Sensor Network Programming Abstractions
 - Represent the physical world to the programmer
 - Abstract distributed interaction with the physical environment





Challenge 3: Distributed Group Management Theory

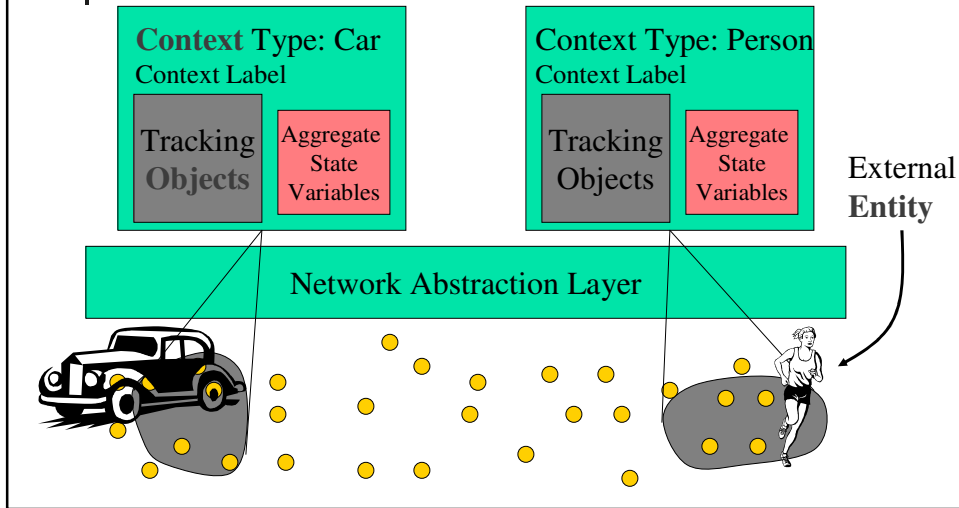
- n Individual nodes or events are insignificant – groups and aggregate state are good abstractions
- n Leverage group communication (a mature topic)
 - n Real-time
 - n Temporal constraints à spatiotemporal constraints
- n New group types/semantics:
 - n Groups may be very dynamic, high failure-rate, variable topology, migration/mobility, relaxed semantics
 - n Semantics cover interaction with the environment
 - n Different interactions give rise to different group types
 - n New semantics, group properties and invariants
 - n Type-specific group communication protocols
 - n Programming interface and primitives
- n Example: Environmentally Immersive Programming



Environmentally Immersive Programming

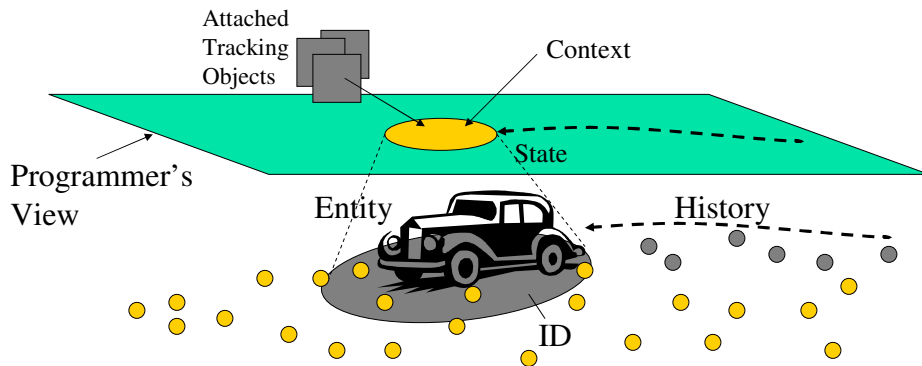
- n Exports a new address space in which the addressed entities (called *contexts*) are representations of physical entities in the external environment
- n *Contexts*:
 - n Logical representations of entities in the external world
 - n Have unique names (context labels) – same as IP hosts
 - n Instantiated when the corresponding external entities are observed in the environment – follow these entities around
 - n Tracking objects (tasks) can be attached to contexts to execute in the vicinity of the corresponding real-world entity
- n Tasks (attached to contexts) can communicate and invoke each other's methods remotely

Programming Model



Contexts and Objects

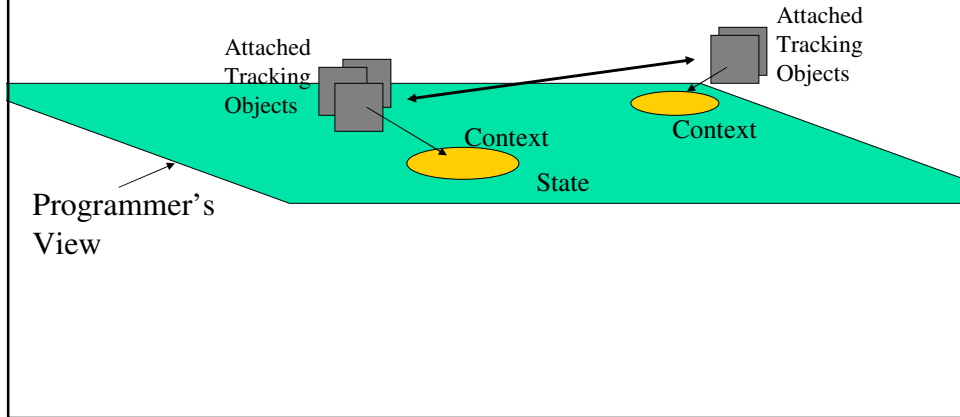
- n Contexts: Encapsulate entity state and tracking objects
- n Tracking objects: Perform entity-specific computation, communication and sensing





Communication

n Objects may export methods for remote invocation

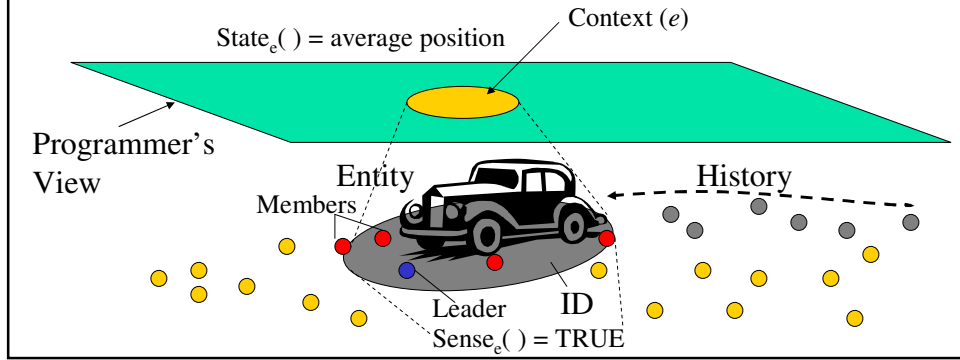


Context Example

```

begin context tracker
  sense: magnetic() + motion();
  state: location = avg (position);
end

```



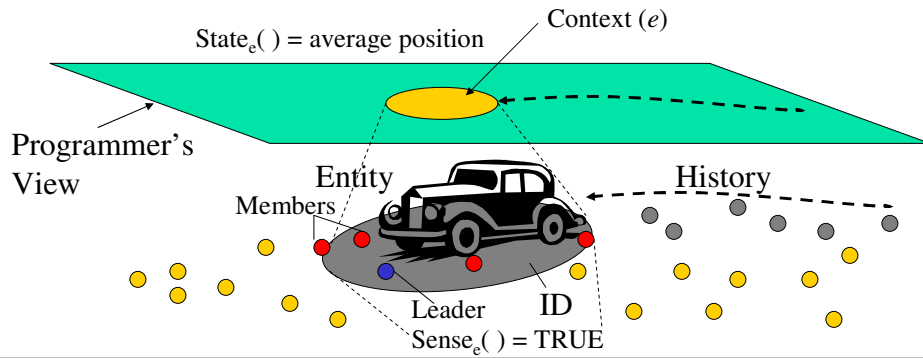
Attaching Objects

```

begin context tracker
  sense: magnetic() + motion();
  state: location = avg (position);
end
  
```

```

begin object reporter
  send (state, home);
end
  
```



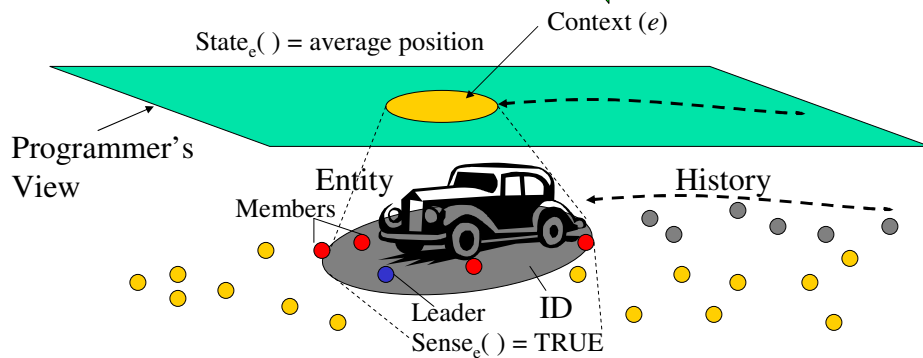
Attaching Objects

```

begin context tracker
  sense: magnetic() + motion();
  state: location = avg (position);
end
  
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```

begin object reporter
  send (state, home);
end
  
```



Attaching Objects

```

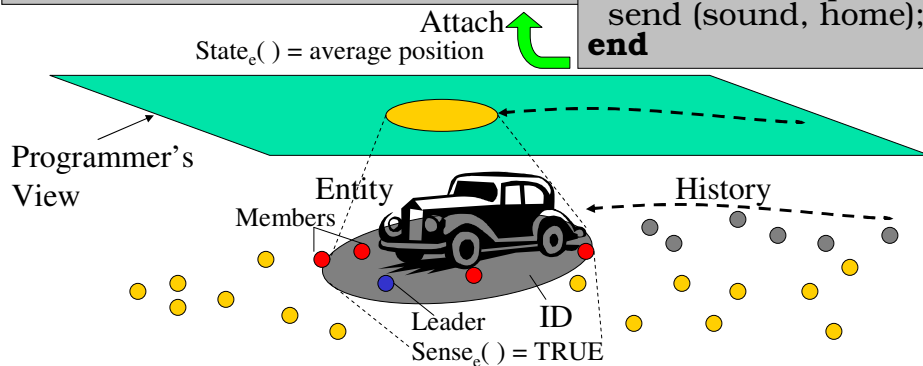
begin context tracker
  sense: magnetic() + motion();
  state: location = avg (position);
end
  
```

```

begin object reporter
  send (state, home);
end
  
```

```

begin object mic
  turn-on microphone;
  send (sound, home);
end
  
```




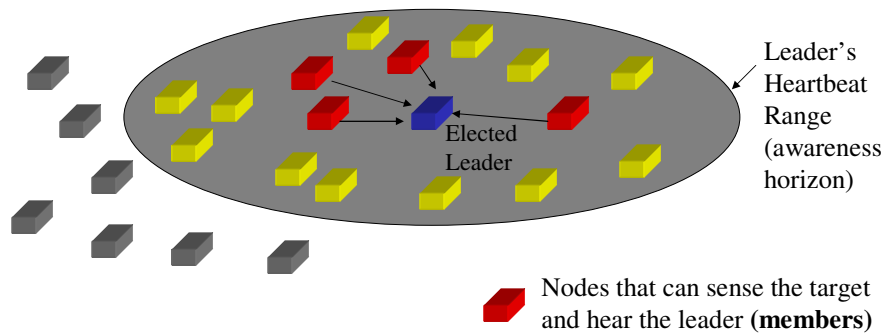
Challenges

- n Context maintenance
 - n Context definition and instantiation (discovery of external entities)
 - n Aggregate state management (abstract state of a dynamically changing group of sensor nodes)
- n Unique context representation
 - n Uniqueness
 - n Context migration and handoff (track locales of mobile physical entities in the environment)
- n Communication, entity location and remote invocation

Group Management and Context Formation

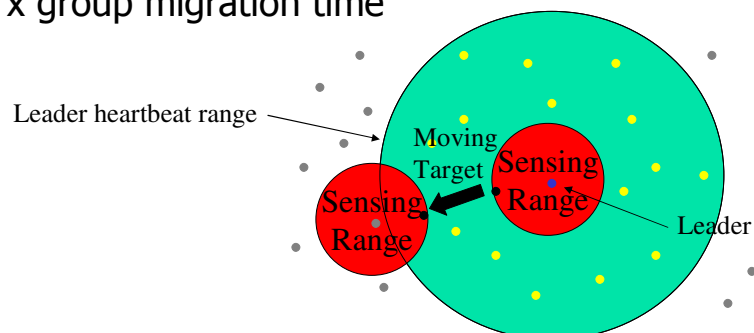
- n Nodes sensing a given target form a single context
- n Context has leader, members, **and followers**

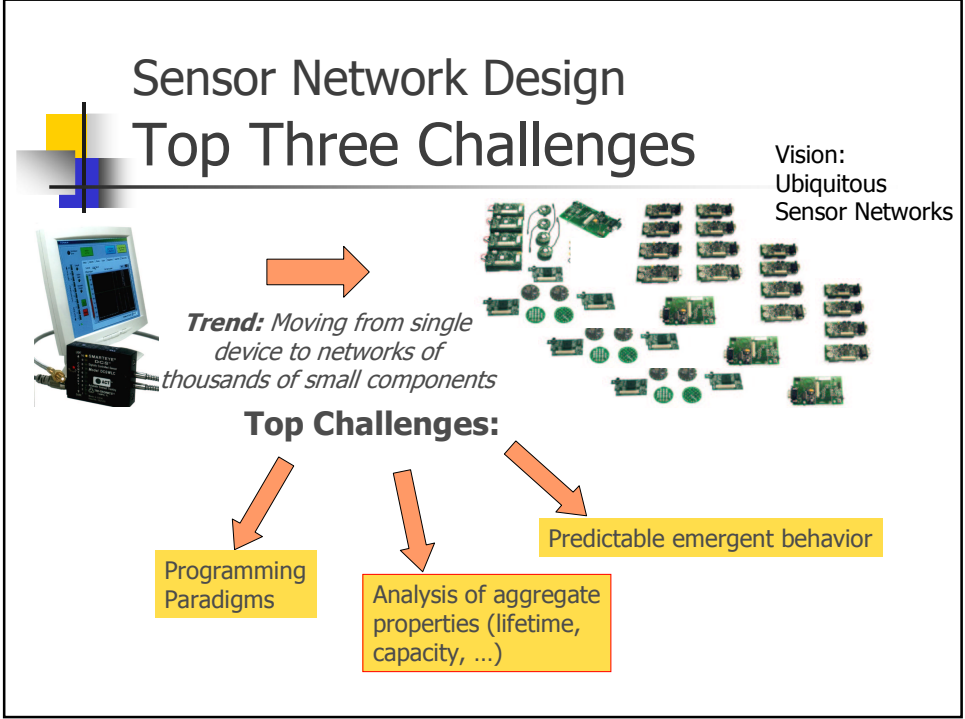
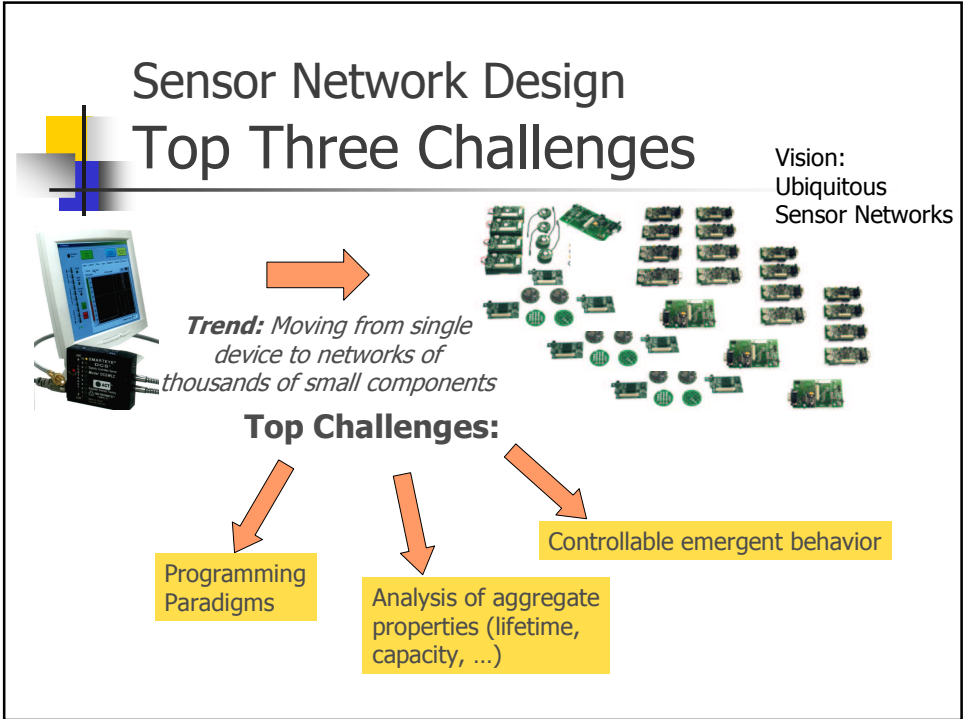
 Nodes that cannot sense the target but hear the leader (**followers**)



A New Spatiotemporal Problem: Unique Representation

- n Target can't move fast enough to be sensed by a node that is outside leader heartbeat range:
- n $\text{Communication range} > 2 \text{ Sensing Range} + \text{Target speed} \times \text{group migration time}$





Real-Time Analysis of Sensor Networks

- n What is the relation between radio range, network density, total number of nodes, number of sinks, packet length, packet scheduling policy, MAC-layer protocol, and end-to-end **packet ability to meet deadlines?**
- n How to choose network parameters that satisfy delay bounds on communication?
- n Optimality (of MAC scheduling policies), capacity proofs, convergence proofs, etc.

Real-Time Objects

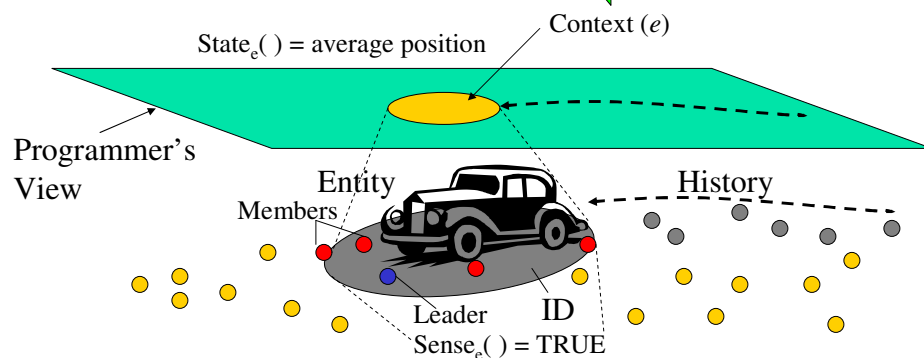
```

begin context tracker
  sense: magnetic() + motion();
  state: location = avg (position,3,2);
end
  
```

```

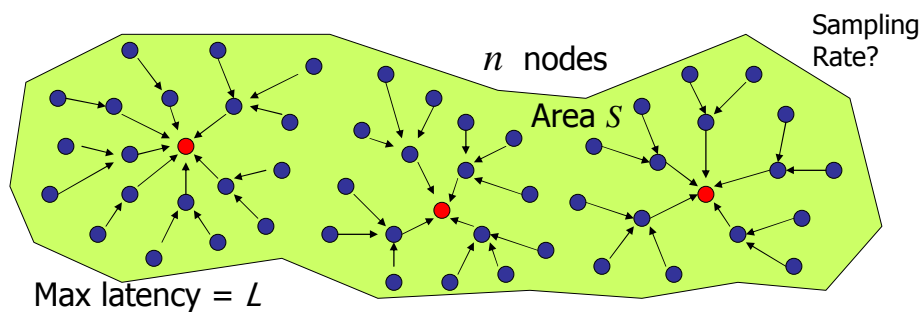
begin RObject reporter
  send (state, home,30s);
end
  
```

Attach



Capacity Planning for Real-Time Wireless Sensor Networks

- n Seminal recent work established wireless network capacity bounds
- n What if traffic has deadlines and only bits that make it by the latency constraint are counted towards throughput?
- n Problem: express *real-time* network capacity that quantifies the throughput of timely bits only as a function of network parameters and time constraints



Network Real-time Capacity

- n Network bandwidth is the bottleneck (communication scheduling problems)
 - n Task processing time \rightarrow packet transmission time
 - n Scheduler queue \rightarrow network queue
- n Intuitively, network schedulability *decreases* with:
 - n Increased packet size (task processing time), C
 - n Increased distance between source and destination, L
 - n Decreased end-to-end latency constraint, D
- n Schedulability decreases with CL/D
- n Is there a bound $Capacity_{RT}$ such that all packets, i , reach their destinations by their deadlines if:

$$\sum_i \frac{C_i L_i}{D_i} \leq Capacity_{RT}$$

The Stage Delay Theorem

- Let us define synthetic utilization at a resource as:
 $\sum C_i/D_i$ for all eligible tasks through that resource (those that have arrived to the system but whose end-to-end deadline has not expired)

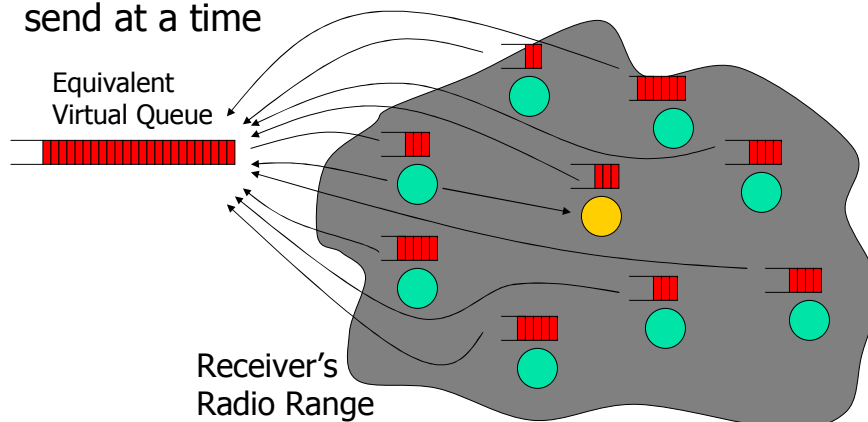
The Stage Delay Theorem: If the synthetic utilization of resource j , does not exceed U_j , then no task is queued on resource j for more than a fraction β_j of its end-to-end deadline*, where:


$$\beta_j = U_j (1 - U_j/2)/(1 - U_j)$$

*This is under deadline monotonic scheduling. Similar results derived for other policies.

The Single Hop Problem

- Only one node in the vicinity of a receiver can send at a time





Throughput Optimization: Maximizing Real-Time Capacity

n Consider a localized communication pattern where each node communicates with nodes at most N hops away.

n On any path, maximize $\sum_j U_j$ subject to

$$\sum_{j=1}^N \frac{U_j(1-U_j/2)}{1-U_j} \leq 1$$

n From symmetry, $U_j = U$

$$\frac{U(1-U/2)}{1-U} = 1/N$$

n Hence, $U = \left(1 + \frac{1}{N} - \sqrt{1 + (\frac{1}{N})^2}\right)$



Real-Time Capacity

n **The total capacity theorem:** In a **load-balanced** network of n nodes, each with a radio of transmission speed W and m neighbors on average, if communication is localized within at most N hops:

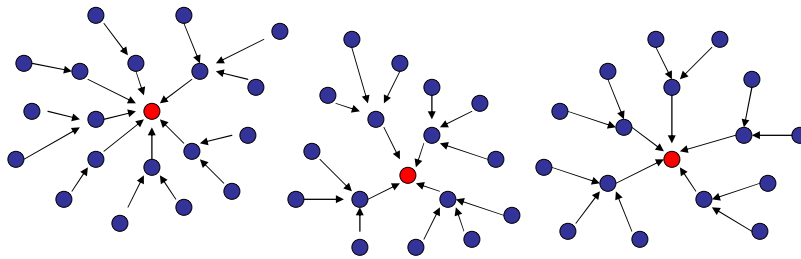
$$Capacity_{opt} = \frac{nW}{m} \left(1 + \frac{1}{N} - \sqrt{1 + (\frac{1}{N})^2}\right)$$

n For large N :

$$Capacity_{opt} \approx \frac{nW}{mN}$$

Real-time Capacity of Multi-hop Data-Collection in Sensor Networks

- n In a **data collection** sensor network with K collection points, maximum path length N , and radio transmission speed W , what is a sufficient bound on real-time capacity?

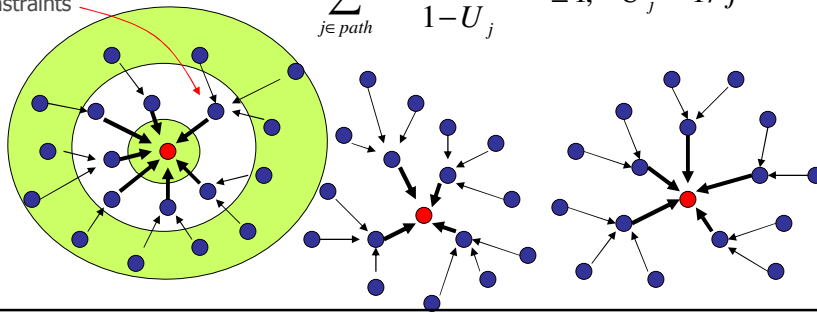


Real-time Capacity of Multi-hop Data-Collection in Sensor Networks

- n In a **data collection** sensor network with K collection points, maximum path length N , and radio transmission speed W , what is a sufficient bound on real-time capacity?

Data conservation constraints

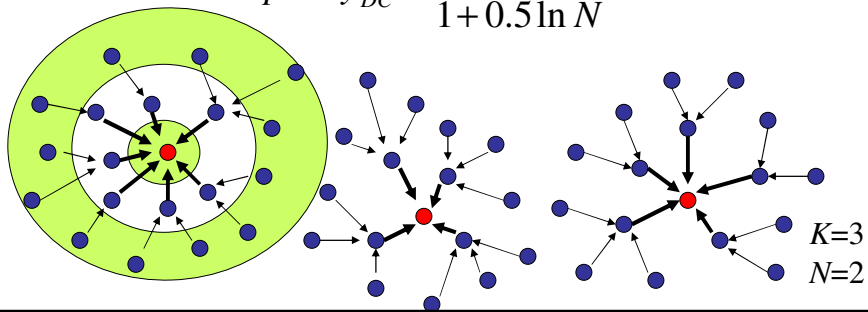
$$\sum_{j \in \text{path}} \frac{U_j(1-U_j/2)}{1-U_j} \leq 1, \quad U_j \propto 1/j$$



Real-time Capacity of Multi-hop Data-Collection in Sensor Networks

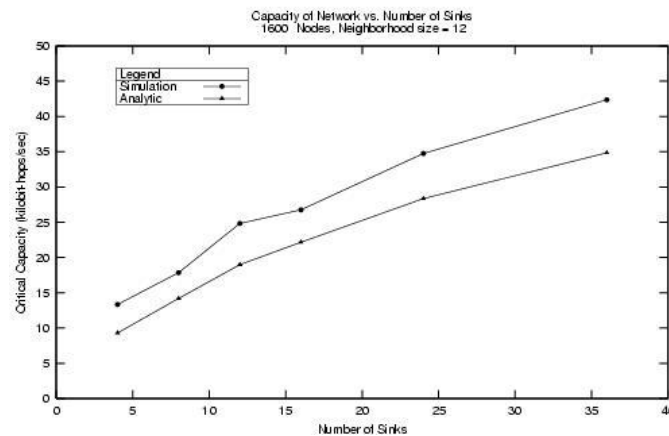
- In a **data collection** sensor network with K collection points, maximum path length N , and radio transmission speed W , a sufficient bound on real-time capacity is:

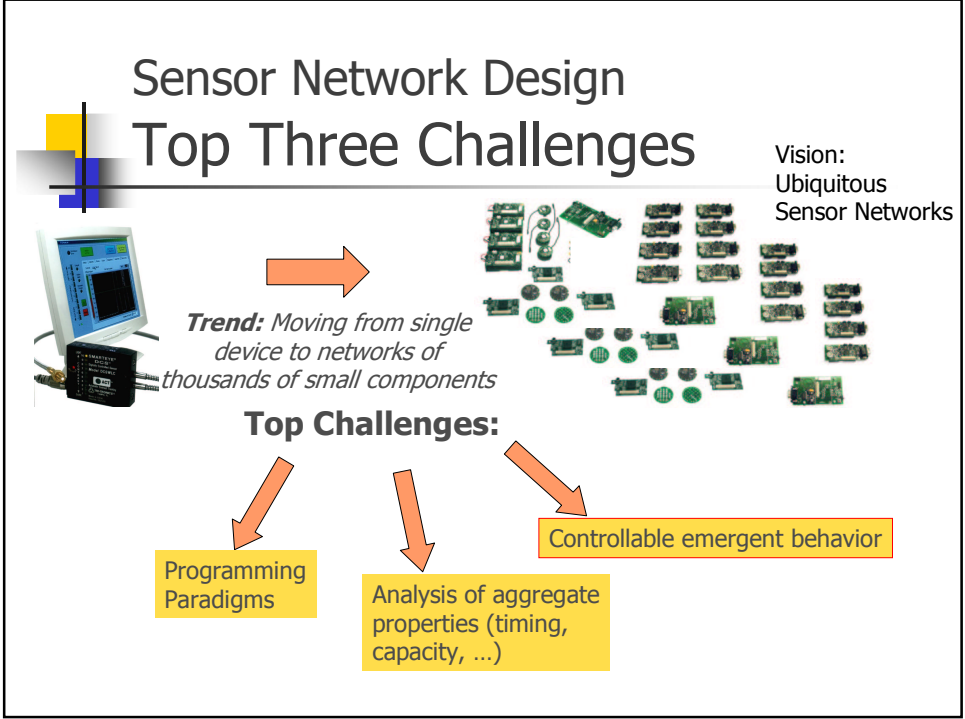
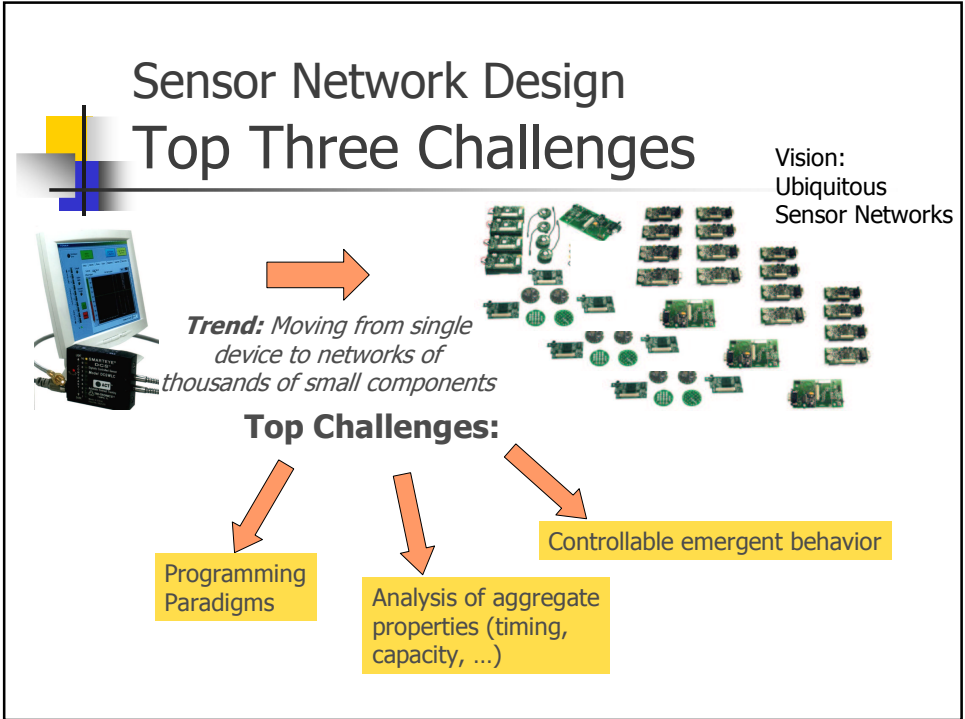
$$Capacity_{DC} \approx \frac{KNW}{1 + 0.5 \ln N}$$

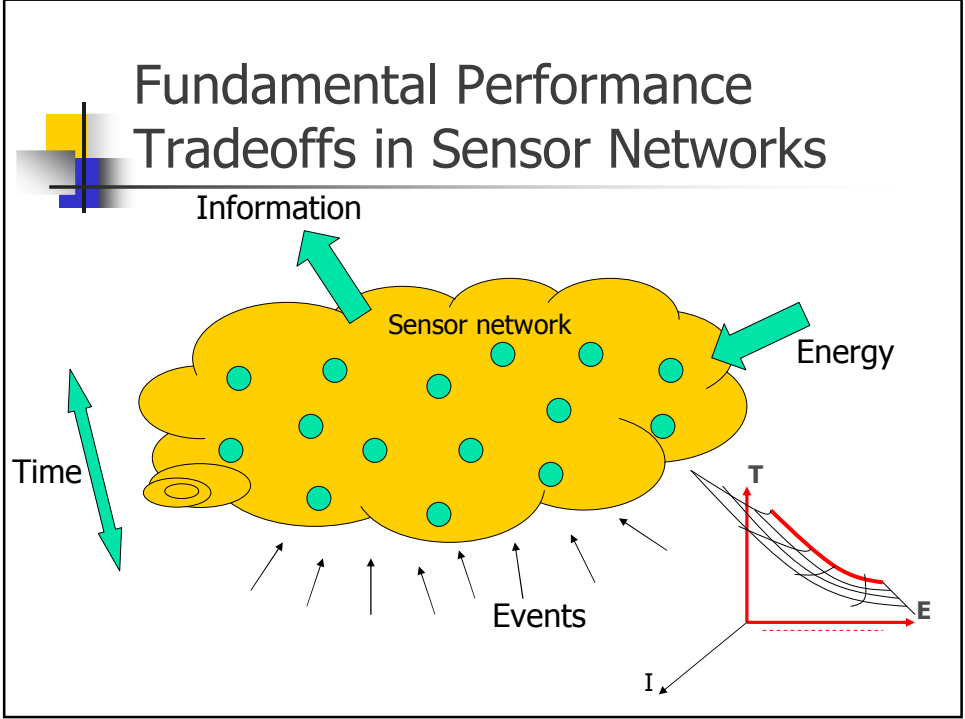
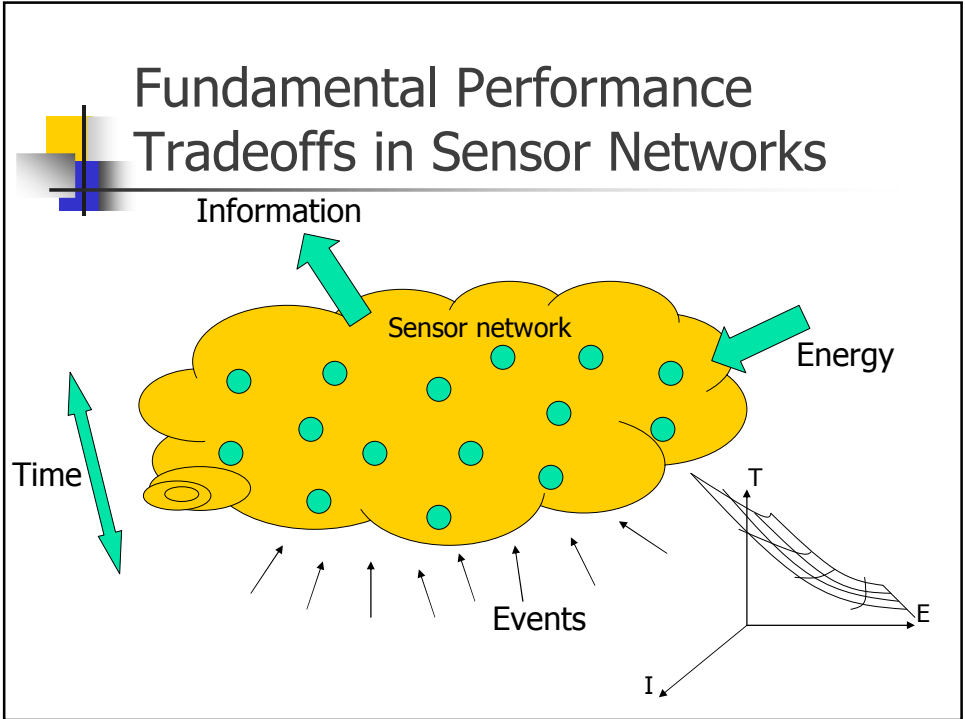


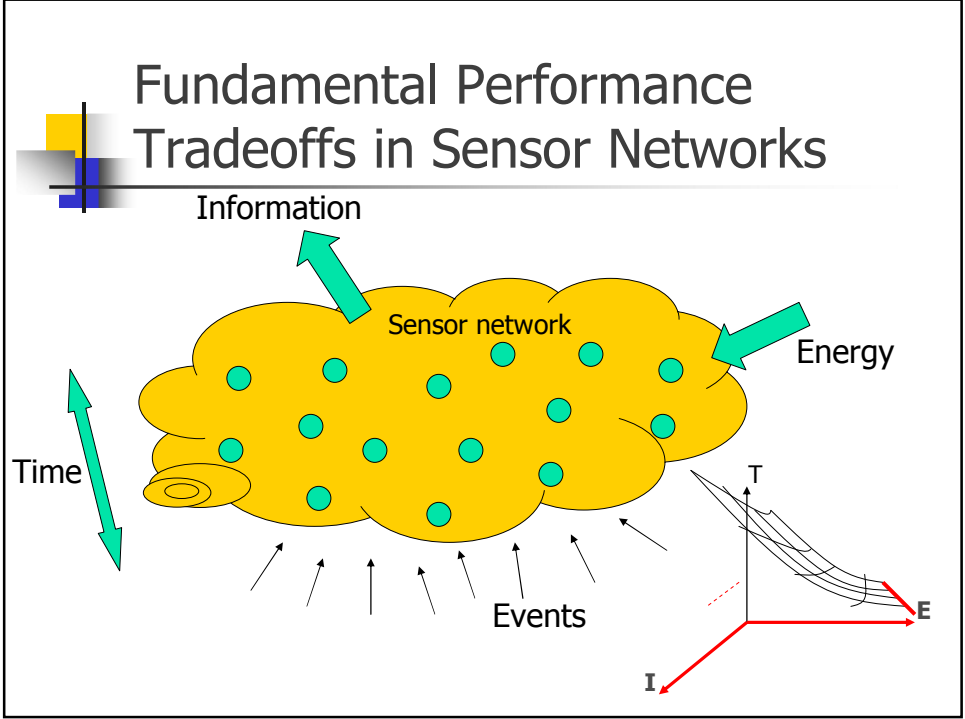
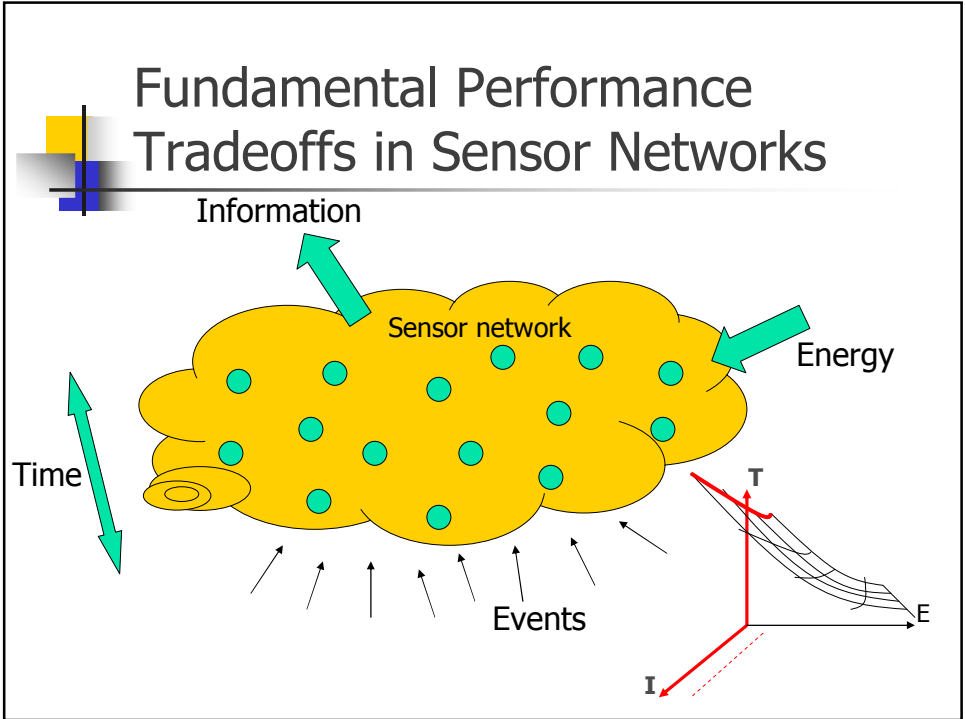
Evaluation: How Pessimistic is Real-time Capacity?

- Simulation versus analytic prediction of the onset of deadline misses in a 1600 node network

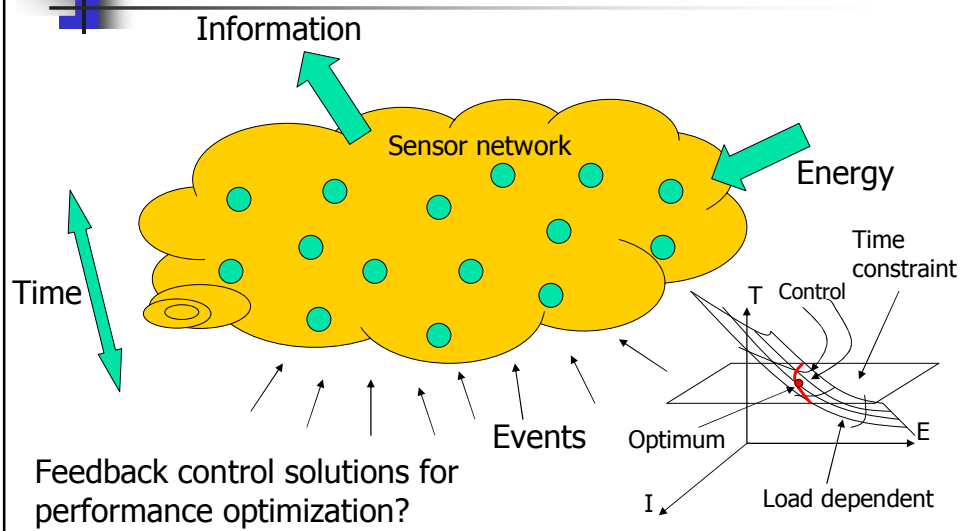








Constrained Optimization and Feedback Control in Sensor Networks



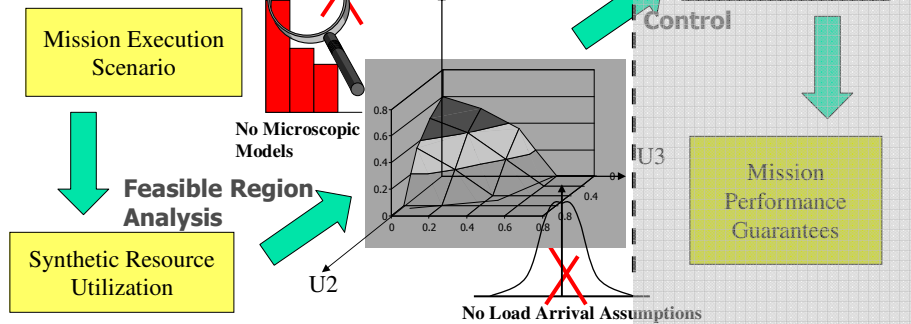
Convergence of Global Behavior in Sensor Networks

- n Sensor network protocols are localized; nodes act independently, locally, in response to local stimuli.
- n How to argue about the global effects of such localized protocols when performed by all nodes?
- n How to induce and analyze convergence to desired global properties?
- n Inspirations from control theory (convergence, stability), Markov decision theory (Markov chains, stochastic models), biology (behavior of social insects, swarm intelligence, bio-differentiation), physics (phase transitions, crystallization), ...

Designing Adaptive Performance-Assured Sensor Networks

- n Achieving desired temporal behavior without fine-grained knowledge
 - n Compute feasibility boundaries
 - n Control the system not to escape them

1. Offline analysis



Final Word: **Sufficient Simplicity**

An Underlying Design Principle

Envisioned solution
Techniques should be:

- **deterministic** analysis
- **aperiodic** tasks
- very **simple** analysis
- **high-level** models
- only **sufficient**

Less error prone (simplicity),
more scalable (simplicity),
and safe (sufficiency)

Trend

Increasing scale/complexity
Decreasing resource cost
Increasing degree of interaction
Less predictable dynamics

Current solution
techniques are:

- exact
- complex
- resource
- optimal



Conclusions

- n Embedded computing is of growing importance in computer science as a discipline
- n New computing paradigms and abstractions are needed
- n Real-time analysis must catch up with sensor network realities
 - n Fundamentally new problems with space and time constraints
 - n Aggregate results as opposed to microscopic models
- n Convergence and aggregate behavior analysis is to be developed
 - n Real-time control of aggregate network properties
- n The science of sensor network design is yet to emerge